## EGERVÁRY RESEARCH GROUP ON COMBINATORIAL OPTIMIZATION



TECHNICAL REPORTS

TR-2010-07. Published by the Egerváry Research Group, Pázmány P. sétány 1/C, H-1117, Budapest, Hungary. Web site: www.cs.elte.hu/egres. ISSN 1587-4451.

# Highly connected rigidity matroids have unique underlying graphs

Tibor Jordán and Viktória E. Kaszanitzky

June 27, 2010

# Highly connected rigidity matroids have unique underlying graphs

Tibor Jordán<sup>\*</sup> and Viktória E. Kaszanitzky<sup>\*\*</sup>

#### Abstract

In this note we consider the following problem: is there a (smallest) integer  $k_d$  such that every graph G is uniquely determined by its d-dimensional rigidity matroid  $\mathcal{R}_d(G)$ , provided that  $\mathcal{R}_d(G)$  is  $k_d$ -connected? Since  $\mathcal{R}_1(G)$  is isomorphic to the cycle matroid of G, a celebrated result of H. Whitney implies that  $k_1 = 3$ . We prove that if G is 7-vertex-connected then it is uniquely determined by  $\mathcal{R}_2(G)$ . We use this result to deduce that  $k_2 \leq 11$ , which gives an affirmative answer for d = 2.

### 1 Introduction

Let  $\mathcal{M}$  be a matroid on ground set E with rank function r and let k be a positive integer. We say that a partition (X, Y) of E is a *k*-separation if

 $\min\{|X|, |Y|\} \ge k$ , and  $r(X) + r(Y) \le r(E) + k - 1.$ 

The connectivity of  $\mathcal{M}$ , denoted by  $\lambda(\mathcal{M})$ , is defined to be the smallest integer j for which  $\mathcal{M}$  has a j-separation. Note that  $\lambda(\mathcal{M}) \geq 1$  for all matroids  $\mathcal{M}$ . We say that  $\mathcal{M}$  is *h*-connected if  $\lambda(\mathcal{M}) \geq h$  holds. We refer the reader to [7] for more details on matroids and matroid connectivity.

The following problem was recently proposed by Brigitte and Herman Servatius [1, Problem 17].

**Problem** Let G be a graph and  $\mathcal{R}_d(G)$  its d-dimensional generic rigidity matroid. Is there a (smallest) constant  $k_d$  such that G is uniquely determined by  $\mathcal{R}_d(G)$  provided that  $\mathcal{R}_d(G)$  is  $k_d$ -connected?

<sup>\*</sup>Department of Operations Research, Eötvös University, Pázmány sétány 1/C, 1117 Budapest, Hungary. Supported by the MTA-ELTE Egerváry Research Group on Combinatorial Optimization and the Hungarian Scientific Research Fund grant no. K81472. e-mail: jordan@cs.elte.hu

<sup>\*\*</sup>Department of Operations Research, Eötvös University, Pázmány sétány 1/C, 1117 Budapest, Hungary.

The d-dimensional generic rigidity matroid (or simply rigidity matroid)  $\mathcal{R}_d(G)$  of graph G = (V, E) is defined on the edge set of G, see [2, 8]. It is not hard to see that  $\mathcal{R}_1(G)$  is isomorphic to the cycle matroid of G, which implies, by a theorem of H. Whitney [9], that  $k_1 = 3$ . The two-dimensional rigidity matroid was characterized by G. Laman [5], who proved that a set  $F \subseteq E$  is independent in  $\mathcal{R}_2(G)$  if and only if

$$|E(G[X]) \cap F| \le 2|X| - 3 \quad \text{for all} \quad X \subseteq V \quad \text{with} \quad |X| \ge 2, \tag{1}$$

where G[X] denotes the subgraph of G induced by X. Note that  $r_2(E) \leq 2|V| - 3$  by (1). In this note we shall prove that  $k_2$  exists and provide an explicit bound  $k_2 \leq 11$ . It is a major open problem to find a good characterization for independence in d-dimensional rigidity matroids, for  $d \geq 3$ . Thus the problem for higher dimensions is probably substantially harder.

We shall consider graphs without loops and isolated vertices. Henceforth we shall assume that d = 2 and omit the subscripts referring to the dimension.

#### 2 Highly connected graphs

We first show that if G is highly connected then its rigidity matroid uniquely determines G. We need some more definitions. Let G = (V, E) be a graph. We say that G is rigid if r(E) = 2|V| - 3 and that G is redundantly rigid if G - e is rigid for all  $e \in E$ . A k-vertex separation of a graph H = (V, E) is a pair  $(H_1, H_2)$  of edgedisjoint subgraphs of G each with at least k + 1 vertices such that  $H = H_1 \cup H_2$  and  $|V(H_1) \cap V(H_2)| = k$ . The graph is said to be k-vertex-connected if it has at least k + 1 vertices and has no j-vertex separation for all  $0 \le j \le k - 1$ .

We shall also need the following three results from combinatorial rigidity.

**Lemma 2.1.** [2, Theorem 4.7.2], [4, Lemma 3.1] Suppose that  $\mathcal{R}(G)$  is 2-connected. Then G is redundantly rigid.

Theorem 2.2. [6, Theorem 2] Every 6-vertex-connected graph is redundantly rigid.

**Theorem 2.3.** [4, Theorem 3.2] Suppose that G is 3-vertex-connected and redundantly rigid. Then  $\mathcal{R}(G)$  is 2-connected.

The proof method of our first result is motivated by a proof for (a special case of) Whitney's theorem, due to J. Edmonds (see [7]). Let  $J \subseteq E$  be a set of elements in matroid  $\mathcal{M}$ . We say that J is a 2-hyperplane of  $\mathcal{M}$  if r(J) = r(E) - 2 and for all  $e \in E - J$  we have r(J + e) = r(E) - 1.

**Theorem 2.4.** Let G and H be two graphs and suppose that  $\mathcal{R}(G)$  is isomorphic to  $\mathcal{R}(H)$ . If G is 7-vertex-connected then G is isomorphic to H.

*Proof.* We say that a 2-hyperplane J of  $\mathcal{R}(G)$  is 2-connected if the matroid restriction of  $\mathcal{R}(G)$  to J is 2-connected. Since G is 7-vertex-connected, Theorems 2.2 and 2.3 imply that G is rigid and E(G - v) (i.e. the edge set E minus the vertex bond of v) is a 2-connected 2-hyperplane of  $\mathcal{R}(G)$  for all  $v \in V(G)$ . Now consider an arbitrary 2-connected 2-hyperplane J of  $\mathcal{R}(G)$ . By Lemma 2.1 the subgraph L = (V(J), J) of G on the set of end vertices of J is rigid. Thus r(J) = 2|V(J)| - 3 and, since 2-hyperplanes are closed sets, it follows that L is an induced subgraph of G. By using the fact that G is rigid, we obtain |V(G)| = |V(J)| + 1. Thus the complement of J corresponds to a vertex bond of G.

It follows that there is a bijection between V(G) and the 2-connected 2-hyperplanes of  $\mathcal{R}(G)$  and that  $\mathcal{R}(G)$  uniquely determines the vertex-edge incidencies in G.

By the assumption of the theorem  $\mathcal{R}(G)$  and  $\mathcal{R}(H)$  are isomorphic. It follows from Theorems 2.2 and 2.3 that  $\mathcal{R}(G)$  is 2-connected. Thus  $\mathcal{R}(H)$  is also 2-connected and hence H is rigid by Lemma 2.1. This implies that 2|V(G)| - 3 = r(G) = r(H) =2|V(H)| - 3 and hence |V(G)| = |V(H)|. Thus  $\mathcal{R}(H)$  has |V(H)| 2-connected 2hyperplanes. So G and H are isomorphic, as claimed.

#### 2.1 Examples

The bound on the connectivity of G in Theorem 2.4 might be improved to 6, but it cannot be replaced by 5. To prove this claim we recall the following family of graphs from [6]: let G be a 5-regular 5-vertex-connected graph on k vertices. Split every vertex of G into 5 vertices of degree one, and identify these 5 vertices with the vertices of a complete graph  $K_5$  on 5 vertices. See Figure 1 for two (non-isomorphic) examples with k = 8.

It is easy to see that the resulting graph G' on 5k vertices is 5-vertex-connected. It is also easy to verify that G' has rank at most  $\frac{19}{2}k$ , hence G' is not rigid when  $k \geq 8$ , see [6]. Furthermore, by using the Henneberg inductive construction to verify independence, one can also show that the rank of G' is exactly  $\frac{19}{2}k$  and that the deletion of an arbitrary edge connecting distinct  $K_5$ 's decreases the rank by one. Thus  $\mathcal{R}(G')$  is the direct sum of k copies of  $\mathcal{R}(K_5)$  and  $\frac{5}{2}k$  copies of  $\mathcal{R}(K_2)$ , for any choice of the initial graph G. Our claim follows, since there exist non-isomorphic 5-regular 5-vertex-connected graphs on  $k \geq 8$  vertices for all  $k \geq 8$ .

We also have similar examples with rigid graphs, but with smaller connectivity. The graphs on Figure 2 are non-isomorphic 3-vertex-connected rigid graphs of the same size. Their rigidity matroids are isomorphic, since the edge set of both graphs is a circuit in the corresponding rigidity matroid. This implies that 7-vertex-connected cannot be replaced by 3-vertex-connected in Theorem 2.4, even if we add the assumption that G is rigid.

#### 3 Highly connected matroids

In this section we show that highly connected rigidity matroids have unique underlying graphs. We shall need the following two lemmas and Theorem 2.4. Let d(v)denote the degree of vertex v in G and let  $\delta(G) = \min\{d(v) : v \in V(G)\}$  denote the minimum degree of G.

**Lemma 3.1.** Let G = (V, E) be a rigid graph on at least three vertices and suppose that  $\mathcal{R}(G)$  is k-connected for some  $k \geq 1$ . Then  $\delta(G) \geq k + 1$ .

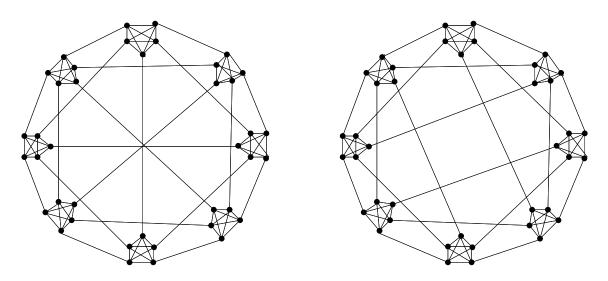


Figure 1: Two non-isomorphic 5-vertex-connected graphs with isomorphic rigidity matroids.

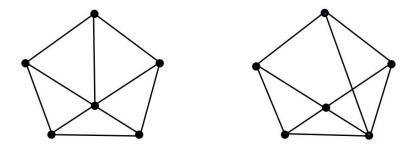


Figure 2: Two non-isomorphic 3-vertex-connected rigid graphs whose edge sets are circuits in their rigidity matroids.

Proof. Since G is rigid, G is 2-vertex connected and  $\delta(G) \geq 2$ . Let X be the set of edges obtained from the vertex bond of some vertex v of degree d(v) by deleting an arbitrary edge. Let Y = E - X. The 2-vertex connectivity of G implies that  $|Y| = |E(G - v)| + 1 \geq |V(G)| - 1 + 1 \geq d(v)$ . Thus  $\min\{|X|, |Y|\} \geq d(v) - 1$  holds. Since X is a co-circuit of  $\mathcal{R}(G)$ , we have

$$r(X) + r(Y) \le d(v) - 1 + r(E) - 1 = r(E) + d(v) - 2.$$

Hence (X, Y) is a (d(v) - 1)-separator of  $\mathcal{R}(G)$ , which implies  $\delta(G) \geq k + 1$ , as required.

**Lemma 3.2.** Let G = (V, E) be a graph and suppose that  $\mathcal{R}(G)$  is (2k-3)-connected for some  $k \geq 3$ . Then G is k-vertex connected.

*Proof.* The hypothesis of the lemma implies that  $\mathcal{R}(G)$  is 2-connected. Thus G is rigid by Lemma 2.1. Hence r(E) = 2|V| - 3 and, by Lemma 3.1, we have  $\delta(G) \ge 2k - 2$  and  $|V| \ge 2k - 1 \ge k + 1$ .

For a contradiction suppose that G has a *j*-vertex separation  $(G_1, G_2)$  for some  $j \leq k-1$ . Let  $X = E(G_1)$  and  $Y = E(G_2)$ . Since  $\delta(G) \geq 2k-2$ , we must have  $\min\{|X|, |Y|\} \geq 2k-2$ . By using (1) we can now deduce that

$$r(X) + r(Y) \le 2|V(G_1)| - 3 + 2|V(G_2)| - 3 = 2(|V| + j) - 6 \le 2|V| + 2k - 8 = r(E) + 2k - 5.$$

Hence (X, Y) is a (2k-4)-separator of  $\mathcal{R}(G)$ , a contradiction. This proves the lemma.

Note that a highly vertex-connected graph G does not necessarily have a highly connected rigidity matroid. The existence of a complete graph  $K_4$  in G (whose edge set is a circuit in  $\mathcal{R}(G)$ ) implies that  $\lambda(\mathcal{R}(G)) \leq 6$ , even if G is highly vertex-connected.

The main result of this section is now a direct corollary of Theorem 2.4 and Lemma 3.2.

**Theorem 3.3.** Let G and H be two graphs and suppose that  $\mathcal{R}(G)$  is isomorphic to  $\mathcal{R}(H)$ . If  $\mathcal{R}(G)$  is 11-connected then G is isomorphic to H.

Theorem 3.3 implies that  $k_2 \leq 11$ . By the example of Figure 2 we have  $k_2 \geq 3$ .

We remark that the proofs and results in this section can easily be extended to *vertical connectivity*, which is another natural form of matroid connectivity [7]. In particular, we can replace 11-connected by vertically 11-connected in Theorem 3.3.

#### 4 Concluding remarks

In this note we have shown that a highly connected two-dimensional rigidity matroid uniquely determines its underlying graph. Since no good characterization is known for independence in the three-dimensional rigidity matroid, the question whether  $k_3$ exists seems more difficult. We note that three-dimensional versions of some of the key results that we used in the proofs exist as conjectures: Lovász and Yemini [6] conjecture that 12-vertex-connected graphs are rigid in three-space, while Jackson and Jordán [3] conjecture that if G is 5-vertex-connected and  $\mathcal{R}_3(G)$  is 2-connected then G is redundantly rigid. The bounds on the vertex connectivity would be best possible in both conjectures.

We thank B. Jackson and B. Servatius for their comments on an earlier version of this paper.

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