# NP-hardness of the Clar number in general plane graphs 

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#### Abstract

We prove that calculating the Clar number in general plane graphs is NPhard.


## 1 Problem Definition

Let $G=(V, E)$ denote a 2-connected planar graph which has a perfect matching. For a planar embedding of $G$ and a perfect matching $M$ of $G$ let $F_{M}$ denote the set of those faces which alternate with respect to $M$. Note that faces in $F_{M}$ are even. A pairwise vertex disjoint subset of $F_{M}$ is a Clar set with respect to $M$. A subset $C$ of the faces is a Clar set if there exists a perfect matching $M$ for which $C$ is a Clar set with respect to $M$. Note that a set of pairwise vertex disjoint even faces is a Clar set if and only if deleting all (the nodes of) these even faces the remaining graph still has a perfect matching. The Clar number of $G$, denoted by $C l(G)$ is the maximum size of a Clar set.

It was proved by Abeledo and Atkinson [T] that the Clar number can be computed in polynomial time if $G$ is bipartite planar.

In this quick proof we show that the general problem is NP-hard.
Theorem 1. It is NP-hard to calculate the Clar number of a planar graph.

## 2 Independent Set Problem

Definition 2. Let $\alpha(G)$ denote the maximum size of an independent set in $G$.
Lemma 3. The Independent Set problem is NP-hard even for planar graphs with odd faces only.

Proof. The independent set problem is NP-hard for planar graphs (see Problem GT20 in [Z] $)$. Let $G=(V, E)$ denote an instance of this problem. If $G$ has an even face $F$, let $G_{F}$ denote the planar graph attained from $G$ by the following operation. We add three vertices $a, b, c$ inside $F$ and edges $a b, b c, c a, a u, b u, b v$ where $u$ and $v$ form an edge of $F$ (see Figure (I).
Claim 4. $\alpha\left(G_{F}\right)=1+\alpha(G)$.


Figure 1: Eliminating even faces.

Proof. First, for an independent set $I$ of $G$, clearly $I \cup c$ is independent in $G_{F}$ and hence $\alpha\left(G_{F}\right) \geq 1+\alpha(G)$. Second, an independent set $I_{F}$ in $G_{F}$ can contain at most one vertex from the set $\{a, b, c\}$. Since $I_{F} \backslash\{a, b, c\}$ is independent in $G$ we get that $\alpha(G) \geq \alpha\left(G_{F}\right)-1$.

Note that the number of even faces of $G_{F}$ is one less than that of $G$. Let $\mathbb{F}$ denote the set of even faces of $G$. By consecutively applying the above operation on every member of $\mathbb{F}$ we get another graph $G_{\mathbb{F}}$ for which $\alpha\left(G_{\mathbb{F}}\right)=\alpha(G)+|\mathbb{F}|$ and which has odd faces only.

## 3 Proof of Theorem

Proof of Theorem [ $\square$. We prove the theorem by reducing the Independent Set problem for planar graphs with odd faces to the Clar number problem. Let $G=(V, E)$ denote a 2-connected instance of this problem (the proof can be easily extended to general connected graphs). We construct graph $G^{\prime}$ the following way: for every edge of $G$ we add two vertices to $G^{\prime}$. Let $u v \in E$ be an edge of $G$ and let $F_{1}$ and $F_{2}$ denote the faces $u v$ is incident to. We add vertices $u v_{F_{1}}$ and $u v_{F_{2}}$ to $G^{\prime}$ along with the edge $u v_{F_{1}} u v_{F_{2}}$. If edges $u v$ and $v w$ are neighboring edges on a face $F$, then we add edge $u v_{F} v w_{F}$ to $G^{\prime}$. It is easy to see that $G^{\prime}$ is planar (see Figure (Z). The faces of $G^{\prime}$ correspond to the faces and vertices of $G$, and if $G$ has odd faces only, then all the even faces of $G^{\prime}$ are the ones corresponding to vertices of $G$. Note that $G^{\prime}$ trivially has a perfect matching $M$ consisting of the edges of the form $u v_{F_{1}} u v_{F_{2}}$, for every $u v \in E$. Since $M$ is alternating on every even face of $G^{\prime}$, corresponding to a vertice of $G$, that is, on every even face of $G^{\prime}$, for this graph the Clar number equals the maximum size of a Clar set with respect to $M$. The Clar sets of $G^{\prime}$ and the independent sets of $G$ have


Figure 2: Reduction of Independent Set
a one to one correspondance, proving the theorem.

Corollary 5. It is also NP-hard to find a maximum cardinality Clar set with respect to a fixed perfect matching.

## 4 Open questions

We proved the NP-hardness of the Clar number for general planar graphs. The classical case of the Clar number, when $G$ has exactly twelve pentagonal faces and every other face is a hexagon, however, is still open. If we were able to specialize the Independent Set problem further to 3-regular planar graphs with odd faces, we would get that the Clar number is NP-hard for graphs whith only hexagonal even faces.

## References

[1] H. G. Abeledo and G. W. Atkinson, A min-max theorem for plane bipartite graphs, Discrete Applied Mathematics 158 (2010), no. 5, 375-378.
[2] M. R. Garey and D. S. Johnson, Computers and intractability: A guide to the theory of NP-completeness, W. H. Freeman \& Co., New York, NY, USA, 1979.

